## Module 1: Proofs

## Operational math bootcamp

Statistical Sciences

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## Outline

- Logic
- Review of Proof Techniques
- Introduction to Set Theory


## Propositional logic

Propositions are statements that could be true or false. They have a corresponding truth value.
ex. " $n$ is odd" and " $n$ is divisible by 2 " are propositions. Let's call them $P$ and $Q$. Whether they are true or not depends on what $n$ is.

We can negate statements: $\neg P$ is the statement " $n$ is not odd"
We can combine statements:

- $P \wedge Q$ is the statement:
- $P \vee Q$ is the statement:

We always assume the inclusive or unless specifically stated otherwise.

## Examples

| Symbol | Meaning |
| :---: | :---: |
| capital letters | propositions |
| $\Longrightarrow$ | implies |
| $\wedge$ | and |
| $\vee$ | inclusive or |
| $\neg$ | not |

- If it's not raining, I won't bring my umbrella.
- I'm a banana or Toronto is in Canada.
- If I pass this exam, I'll be both happy and surprised.


## Truth values

## Example

If it is snowing, then it is cold out.
It is snowing.
Therefore, it is cold out.
Write this using propositional logic:

How do we know if this statement is true or not?

## Truth table

$$
P \Longrightarrow Q
$$

If it is snowing, then it is cold out.

When is this true or false?

| $P$ | $Q$ | $P \Longrightarrow Q$ |
| :---: | :---: | :---: |
| T | T |  |
| T | F |  |
| F | T |  |
| F | F |  |

## Logical equivalence

$$
P \Longrightarrow Q
$$

$$
\neg P \vee Q
$$

| $P$ | $Q$ | $P \Longrightarrow Q$ |
| :---: | :---: | :---: |
| T | T | T |
| T | F | F |
| F | T | T |
| F | F | T |


| $P$ | $Q$ | $\neg P$ | $\neg P \vee Q$ |
| :---: | :---: | :---: | :---: |
| T | T |  |  |
| T | F |  |  |
| F | T |  |  |
| F | F |  |  |

$$
\text { What is } \neg(P \Longrightarrow Q) \text { ? }
$$

## Quantifiers

## For all

"for all", $\forall$, is also called the universal quantifier.
If $P(x)$ is some property that applies to $x$ from some domain, then $\forall x P(x)$ means that the property $P$ holds for every $x$ in the domain.
"Every real number has a non-negative square." We write this as

How do we prove a for all statement?

## Quantifiers

## There exists

"there exists", $\exists$, is also called the existential quantifier.
If $P(x)$ is some property that applies to $x$ from some domain, then $\exists x P(x)$ means that the property $P$ holds for some $x$ in the domain.

4 has a square root in the reals. We write this as

How do we prove a there exists statement?

There is also a special way of writing when there exists a unique element: $\exists$ ! .
For example, we write the statement "there exists a unique positive integer square root of 64" as

## Combining quantifiers

Often we will need to prove statements where we combine quantifiers. Here are some examples:
Statement Logical expression

Every non-zero rational number has a multiplicative inverse

Each integer has a unique additive inverse
$f: \mathbb{R} \rightarrow \mathbb{R}$ is continuous at $x_{0} \in \mathbb{R}$

## Quantifier order \& negation

The order of quantifiers is important! Changing the order changes the meaning. Consider the following example. Which are true? Which are false?

$$
\begin{aligned}
& \forall x \in \mathbb{R} \forall y \in \mathbb{R} x+y=2 \\
& \forall x \in \mathbb{R} \exists y \in \mathbb{R} x+y=2 \\
& \exists x \in \mathbb{R} \forall y \in \mathbb{R} x+y=2 \\
& \exists x \in \mathbb{R} \exists y \in \mathbb{R} x+y=2
\end{aligned}
$$

Negating quantifiers:

$$
\begin{aligned}
\neg \forall x P(x) & =\exists x(\neg P(x)) \\
\neg \exists x P(x) & =\forall x(\neg P(x))
\end{aligned}
$$

The negations of the statements above are:
(Note that we use De Morgan's laws, which are in your exercises:

$$
\neg(P \wedge Q)=\neg P \vee \neg Q \text { and } \neg(P \vee Q)=\neg P \wedge \neg Q .)
$$

Logical expression
Negation
$\forall q \in \mathbb{Q} \backslash\{0\}, \exists s \in \mathbb{Q}$ such that $q s=1$
$\forall x \in \mathbb{Z}, \exists!y \in \mathbb{Z}$ such that $x+y=0$

$$
\begin{aligned}
& \forall \epsilon>0 \exists \delta>0 \text { such that whenever } \mid x- \\
& x_{0}\left|<\delta,\left|f(x)-f\left(x_{0}\right)\right|<\epsilon\right.
\end{aligned}
$$

What do these mean in English?

## Types of proof

- Direct
- Contradiction
- Contrapositive
- Induction


## Direct Proof

Approach: Use the definition and known results.

## Example

## Claim

The product of an even number with another integer is even.

Approach: use the definition of even.

## Direct Proof

## Claim

The product of an even number with another integer is even.

## Definition

We say that an integer $n$ is even if there exists another integer $j$ such that $n=2 j$.
We say that an integer $n$ is odd if there exists another integer $j$ such that $n=2 j+1$.

## Proof.

## Definition

Let $a, b \in \mathbb{Z}$. We say that "a divides $b$ ", written $a \mid b$, if the remainder is zero when $b$ is divided by a, i.e. $\exists j \in \mathbb{Z}$ such that $b=a j$.

## Example

Let $a, b, c \in \mathbb{Z}$ with $a \neq 0$. Prove that if $a \mid b$ and $b \mid c$, then $a \mid c$.

## Proof.

## Claim

If an integer squared is even, then the integer is itself even.

How would you approach this proof?

## Proof by contrapositive

$$
P \Longrightarrow Q
$$

$$
\neg Q \Longrightarrow \neg P
$$

| $P$ | $Q$ | $P \Longrightarrow Q$ |
| :---: | :---: | :---: |
| T | T | T |
| T | F | F |
| F | T | T |
| F | F | T |


| $P$ | $Q$ | $\neg P$ | $\neg Q$ | $\neg Q \Longrightarrow \neg P$ |
| :---: | :---: | :---: | :---: | :---: |
| T | T | F | F |  |
| T | F | F | T |  |
| F | T | T | F |  |
| F | F | T | T |  |

## Proof by contrapositive

## Claim

If an integer squared is even, then the integer is itself even.

## Proof.

## Proof by contradiction

## Claim

The sum of a rational number and an irrational number is irrational.

## Proof.

## Summary

In sum, to prove $P \Longrightarrow Q$ :
Direct proof: assume $P$, prove $Q$
Proof by contrapositive: assume $\neg Q$, prove $\neg P$
Proof by contradiction: assume $P \wedge \neg Q$ and derive something that is impossible

## Induction

## Well-ordering principle for $\mathbb{N}$

Every nonempty set of natural numbers has a least element.

## Principle of mathematical induction

Let $n_{0}$ be a non-negative integer. Suppose $P$ is a property such that
(1) (base case) $P\left(n_{0}\right)$ is true
(2) (induction step) For every integer $k \geq n_{0}$, if $P(k)$ is true, then $P(k+1)$ is true.

Then $P(n)$ is true for every integer $n \geq n_{0}$
Note: Principle of strong mathematical induction: For every integer $k \geq n_{0}$, if $P(n)$ is true for every $n=n_{0}, \ldots, k$, then $P(k+1)$ is true.

## Claim

$$
n!>2^{n} \text { if } n \geq 4(n \in \mathbb{N}) \text {. }
$$

## Proof.

## Claim

Every integer $n \geq 2$ can be written as the product of primes.

## Proof.

We prove this by strong induction on $n$.
Base case:
Inductive hypothesis:
Inductive step:

## References

Gerstein, Larry J. (2012). Introduction to Mathematical Structures and Proofs.
Undergraduate Texts in Mathematics. url: https://link.springer.com/book/10.1007/978-1-4614-4265-3

Lakins, Tamara J. (2016). The Tools of Mathematical Reasoning. Pure and Applied Undergraduate Texts.

